



UNIVERSITÄT  
BAYREUTH

# Descent and Covering Collections

## Part I: Local Solubility

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# Hyperelliptic Curves (1)

Let  $k$  be a field with  $\text{char}(k) \neq 2$ .

A **hyperelliptic curve** over  $k$  is the smooth projective curve associated to an affine plane curve given by an equation of the form

$$y^2 = f(x) = f_n x^n + f_{n-1} x^{n-1} + \dots + f_1 x + f_0$$

with  $f \in k[x]$  **squarefree** (i.e.,  $\text{disc}(f) \neq 0$ ).

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We write

$$C: y^2 = f(x)$$

to denote the **projective** curve  $C$ .

## Hyperelliptic Curves (2)

A more abstract definition (that works over any field  $k$ ) is as follows.

A **hyperelliptic curve** over  $k$  is a **nice** (= smooth, projective and geometrically irreducible) **curve  $C$**  over  $k$  with a map  $\pi: C \rightarrow \mathbb{P}^1$  of degree 2, which is defined over  $k$ .

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Writing  $k(x)$  for the function field of  $\mathbb{P}_k^1$ , the function field of  $C$  is a quadratic extension of  $k(x)$ , so is of the form  $k(x, \sqrt{f(x)})$  if  $\text{char}(k) \neq 2$ . Writing  $y$  for  $\sqrt{f(x)}$ , we obtain the equation  $y^2 = f(x)$ .

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In characteristic 2, one has to consider more general equations of the form

$$y^2 + h(x)y = f(x).$$

## Hyperelliptic Curves (3)

Let  $C: y^2 = f(x)$  be a hyperelliptic curve. Then

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A smooth projective model of  $C$  can be obtained as follows.

Let  $F \in k[x, z]$  be the binary form of degree  $2g + 2$  such that  $F(x, 1) = f(x)$ .

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So the points at infinity on  $C$  are  $\infty_s = (1 : s : 0)$  where  $s^2 = F(1, 0) = f_{2g+2}$ :

There is one point  $\infty = \infty_0$  when  $\deg(f)$  is odd, otherwise there are two (which are  $k$ -rational iff the leading coefficient of  $f$  is a square in  $k$ ).

## Hyperelliptic Curves (4)

Any (nice) curve  $C$  of genus 2 over  $k$  is hyperelliptic over  $k$ :

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Since hyperelliptic curves always have the nontrivial automorphism  $\iota: (x, y) \mapsto (x, -y)$  (called the 'hyperelliptic involution'), there can be curves that are non-isomorphic over  $k$ , but become isomorphic over  $\bar{k}$ .

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In the following, we will concentrate on the case  $k = \mathbb{Q}$   
(or, more generally, an algebraic number field).

The main question will be:

**How can we determine the set  $C(\mathbb{Q})$ ?**

# General Facts on Rational Points on Curves

Recall the following general classification.

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Note that this trichotomy is given by the sign ( $> 0$ ,  $= 0$ ,  $< 0$ ) of the Euler characteristic  $2 - 2g$ , which is a topological invariant of the Riemann surface  $C(\mathbb{C})$ !

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We split the problem into two parts:

- (1) Decide whether  $C(\mathbb{Q})$  is **empty** or not!
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In this course, we will mainly focus on the **first problem**.

## A Simple Test

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Indeed, this proves that  $C(\mathbb{Q}_3) = \emptyset$ !

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Define the **p-adic absolute value** on  $\mathbb{Q}$  by

$$|\xi|_p = \begin{cases} 0 & \text{if } \xi = 0; \\ p^{-n} = p^{-v_p(\xi)} & \text{if } \xi = p^n \frac{a}{b} \text{ with } p \nmid ab. \end{cases}$$

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Then  $|\alpha\beta|_p = |\alpha|_p \cdot |\beta|_p$  and  $|\alpha + \beta|_p \leq \max\{|\alpha|_p, |\beta|_p\} \leq |\alpha|_p + |\beta|_p$ , so we can define  $\mathbb{Q}_p$  as the **completion** of  $\mathbb{Q}$  with respect to  $|\cdot|_p$ .

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The closed unit ball  $\mathbb{Z}_p = \{\xi \in \mathbb{Q}_p : |\xi|_p \leq 1\}$  forms a **compact subring**; it is the topological closure of  $\mathbb{Z}$  in  $\mathbb{Q}_p$ .

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It is then easy to check that  $p^n\mathbb{Z}_p = \{\xi \in \mathbb{Q}_p : |\xi|_p \leq p^{-n}\}$   
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and we can think of elements of  $\mathbb{Q}_p$  as '**Laurent series in p**'

$$\xi = \sum_{n=n_0}^{\infty} a_n p^n \quad \text{with } a_n \in \{0, 1, \dots, p-1\}.$$

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So the following definition makes sense.

## **Definition.**

A (nice) curve  $C$  over  $\mathbb{Q}$  is said to be **everywhere locally soluble** or **ELS**, if  $C(\mathbb{R}) \neq \emptyset$  and  $C(\mathbb{Q}_p) \neq \emptyset$  for all primes  $p$ .

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**Question.** Can we **decide** if a given curve is ELS?

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# Checking for $p$ -adic Points (1)

**Theorem (Weil).**

Let  $C$  be a nice curve of genus  $g$  over  $\mathbb{F}_p$ . Then

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Every point on  $\mathcal{C} \otimes \mathbb{F}_p$  is smooth, so  $C(\mathbb{Q}_p) = \mathcal{C}(\mathbb{Q}_p) \neq \emptyset$  by Hensel. □

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If  $p < 4g^2$  or the curve has **bad reduction at  $p$** ,  
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Since we assume that  **$C$  is smooth**,

the ‘zooming in’ will eventually produce models with smooth fiber over  $\mathbb{F}_p$ .

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So we **can decide** whether  $C$  is ELS or not!